Theory of Estimation-of-Distribution Algorithms

PPSN 2024 Tutorial







- » What are they?
- » Focus on the most commonly theoretically studied ones

2 Theory



(Mathematical rigor)

- » What are they?
- » Focus on the most commonly theoretically studied ones

- » What are they?
- » Focus on the most commonly theoretically studied ones



(Mathematical rigor)

» What is studied?

- » What are they?
- » Focus on the most commonly theoretically studied ones



(Mathematical rigor)

- » What is studied?
- » Selection of important results
 - *** Understand** the basic idea

- » What are they?
- » Focus on the most commonly theoretically studied ones



(Mathematical rigor)

- » What is studied?
- » Selection of important results
 - *** Understand** the basic idea

» Depth over breadth





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X



DOI: 10.1145/3377929.3389888

(more **breadth**)

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DOI: 10.1145/3377929.3389888

(more **breadth**)

Carsten Witt



Sunday, Sep. 15 9:30–12:30

















<i>x</i> ₁	x_2	x_3	x_4	<i>x</i> ₅
0	0	0	0	0
0	0	0	0	1
0	0	0	0	2
		:		



$$\prod_{i \in [\#vars]} (\#vals(x_i)) - 1$$



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compressed, directed



 $x_3 \ x_4 \ x_5$ $x_1 \quad x_2$: .

$$\prod_{i \in [\#vars]} (\#vals(x_i)) - 1$$



compressed, directed

<i>x</i> ₁	<i>x</i> ₂	<i>x</i> ₂	<i>x</i> ₃	
0	0	0	0	
1	1	0	1	
2	2	0	2	
•	•			



:

$$\prod_{i \in [\#vars]} (\#vals(x_i)) - 1$$



 $\sum_{i \in [\#vars]} (\#vals(x_i) - 1) \prod_{y \in pred(x_i)} \#vals(y)$ x_1 x_2 $x_2 \ x_3$ 0 0 0 0 ... : :

compressed, directed



$$\prod_{i \in [\#vars]} (\#vals(x_i)) - 1$$



compressed, directed

$$\sum_{i \in [\#vars]} (\#vals(x_i) - 1) \prod_{y \in pred(x_i)} \#vals(y)$$

Theory

$$x_1$$
 x_2 x_3 \cdots x_n
univariate



$f: \{0,1\}^n \to \mathbf{R}$

» Pseudo-Boolean optimization



$f: \{0,1\}^n \to \mathbf{R}$

- » Pseudo-Boolean optimization
 - *** Global optimum** often 1^n



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» Frequency vector p



$f\!:\!\{0,\!1\}^n\to \mathbf{R}$

- » Pseudo-Boolean optimization
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- » Frequency vector $m{p}$
- » p_i ... probability to sample a 1 at i (green mass)



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Multi-valued analyses



$f\!:\!\{0,\!1\}^n\to \mathbf{R}$

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Multi-valued analyses

» DOI: 10.1016/j.tcs.2024.114622



» Runtime Analysis of a Multi-Valued Compact Genetic Algorithm on Generalized OneMax

PPSN 2024

[Adak, Witt'24]



$f: \{0,1\}^n \to \mathbf{R}$

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- » p_i ... probability to sample a 1 at i (green mass)



- » Run time analysis
 - Number of evaluations of *f* until a global optimum is sampled



Multi-valued analyses

- » DOI: 10.1016/j.tcs.2024.114622
- [Ben Jedidia, Doerr, K.'24]
- » Runtime Analysis of a Multi-Valued Compact Genetic Algorithm on Generalized OneMax



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Multi-valued analyses

» p_i ... probability to sample a 1 at i (green mass)

» DOI: 10.1016/j.tcs.2024.114622



- » Run time analysis
 - Number of evaluations of *f* until a global optimum is sampled



» Favorable albeit not necessary model



» Runtime Analysis of a Multi-Valued Compact Genetic Algorithm on Generalized OneMax



[Adak, Witt'24]

Updating a Frequency (commonly)





Updating a Frequency (commonly)








$$\sum_{\substack{\mathbf{p}_i^{(t+1)} - \mathbf{p}_i^{(t)} \mid \mathbf{p}^{(t)} \\ \text{drift}}} E\left[\mathbf{p}_i^{(t+1)} - \mathbf{p}_i^{(t)} \mid \mathbf{p}^{(t)}\right]$$

7

drift



 $\mathrm{E}\left[\boldsymbol{p}_{i}^{(t+1)} - \boldsymbol{p}_{i}^{(t)} \mid \boldsymbol{p}^{(t)} \right]$ » Estimating the drift, estimates the expected time that the frequency passes a specific value

7



Updating a Frequency (commonly) **Very tricky!**

samples $\left\{ \boldsymbol{x}^{(j)} \right\}_{j \in [\lambda]}$





 $\boldsymbol{p}_i^{(t)}$

 $E\left[\boldsymbol{p}_{i}^{(t+1)} - \boldsymbol{p}_{i}^{(t)} \mid \boldsymbol{p}^{(t)}\right]$ * Estimating the drift, estimates the expected time that the frequency passes a specific value ✤ Drift analysis X 1712.00964 [Lengler'18] drift

(ordered by decreasing fitness)

selected samples $\left\{ \widehat{x}^{(k)} \right\}_{k \in [\mu]}$

X 2406.14589 [Kötzing'24]

» What happens if the samples are not informative?



 $\boldsymbol{p}_i^{(t)}$

» What happens if the samples are not informative?

* Each $\{x_i^{(j)}\}_{j\in[\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k\in[\mu]}$ follows the same distribution

» Can happen regularly during a run



 $\boldsymbol{p}_i^{(t)}$

- » What happens if the samples are not informative?
 - Each $\{x_i^{(j)}\}_{j \in [\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k \in [\mu]}$ follows the same distribution
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Removes the tricky situation

- » What happens if the samples are not informative?
 - Each $\{x_i^{(j)}\}_{j \in [\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k \in [\mu]}$ follows the same distribution
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$$\mathbf{E}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \boldsymbol{p}_{i}^{(t)}$$

balanced

- » What happens if the samples are not informative?
 - Each $\{x_i^{(j)}\}_{j \in [\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k \in [\mu]}$ follows the same distribution
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 Due to random fluctuations; despite a clear signal

Genetic drift

- » What happens if the samples are not informative?
 - Each $\{x_i^{(j)}\}_{j \in [\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k \in [\mu]}$ follows the same distribution
- » Can happen regularly during a run



- Removes the tricky situation $E\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \boldsymbol{p}_{i}^{(t)} \quad \text{* Typical property of univariate EDAs}$ balanced $\boldsymbol{p}_{i} \text{ is a martingale and approaches the borders quickly}$
 - Due to random fluctuations; despite a clear signal

- » What happens if the samples are not informative?
 - Each $\{x_i^{(j)}\}_{j \in [\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k \in [\mu]}$ follows the same distribution
- » Can happen regularly during a run







How fast does p_i cover distance $d \in [0,1]$?

» What happens if the samples are not informative?

Removes the tricky situation

drift via

variance

• Each $\{x_i^{(j)}\}_{j \in [\lambda]}$ and $\{\widehat{x}_i^{(k)}\}_{k \in [\mu]}$ follows the same distribution

 $\mathbb{E}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \boldsymbol{p}_{i}^{(t)}$ » Typical property of univariate EDAs

p_i is a martingale and approaches the borders quickly

» Can happen regularly during a run



- » Due to random fluctuations; despite a clear signal

Genetic drift

balanced



$$\approx \frac{d}{\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right]} \quad \text{[K.'19]}$$

What happens if the samples are **not** informative?

Removes the tricky situation

- Each $\{\boldsymbol{x}_{i}^{(j)}\}_{i \in [\lambda]}$ and $\{\hat{\boldsymbol{x}}_{i}^{(k)}\}_{k \in [\mu]}$ follows the same distribution
- » Can happen regularly during a run





[Doerr, Zheng'20]

Compact Genetic Algorithm (cGA)

[Harik, Lobo, Goldberg'99]

Univariate Marginal Distribution Algorithm (UMDA)

Univariate Marginal Distribution Algorithm (UMDA)



Univariate Marginal Distribution Algorithm (UMDA)



Univariate Marginal Distribution Algorithm (UMDA)



$$\boldsymbol{p}_{i}^{(t+1)} = \boldsymbol{p}_{i}^{(t)} + \frac{1}{K} \left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)} \right)$$

hypothetical population size

Univariate Marginal Distribution Algorithm (UMDA)



$$\boldsymbol{p}_{i}^{(t+1)} = \boldsymbol{p}_{i}^{(t)} + \frac{1}{K} \left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)} \right)$$

hypothetical population size

Univariate Marginal Distribution Algorithm (UMDA)

[Mühlenbein, Paaß'96]



» λ samples each iteration



$$p_i^{(t+1)} = p_i^{(t)} + \frac{1}{K} \left(\hat{x}_i^{(1)} - \hat{x}_i^{(2)} \right)$$

Univariate Marginal Distribution Algorithm (UMDA)



hypothetical population size





$$\boldsymbol{p}_{i}^{(t+1)} = \boldsymbol{p}_{i}^{(t)} + \frac{1}{K} \left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)} \right)$$

hypothetical population size





$$\boldsymbol{p}_i^{(t+1)} = \boldsymbol{p}_i^{(t)} + \frac{1}{K} \left(\widehat{\boldsymbol{x}}_i^{(1)} - \widehat{\boldsymbol{x}}_i^{(2)} \right)$$

hypothetical population size

 $\boldsymbol{p}_i^{(t+1)} = \frac{1}{\mu} \sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_i^{(k)}$



Univariate Marginal Distribution Algorithm (UMDA) (Mühlenbein, Paaß'96) * λ samples each iteration * Select $\mu \le \lambda$ best samples, ranking them * Set to the relative number of 1s

 $\boldsymbol{p}_i^{(t+1)} = \frac{1}{u} \sum_{k \in [u]} \widehat{\boldsymbol{x}}_i^{(k)}$

$$\boldsymbol{p}_{i}^{(t+1)} = \boldsymbol{p}_{i}^{(t)} + \frac{1}{K} \left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)} \right)$$
hypothetical population size



» Population-based incremental learning [Baluja'94]
 » 2-Max-Min ant system with iteration-best update [Neumann, Sudholt, Witt'10]





$$\boldsymbol{p}_{i}^{(t+1)} = \boldsymbol{p}_{i}^{(t)} + \frac{1}{K} \left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)} \right)$$
hypothetical population size

 $\boldsymbol{p}_i^{(t+1)} = \frac{1}{\mu} \sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_i^{(k)}$

Population-based incremental learning

[Baluja'**94**]

» 2-Max-Min ant system with iteration-best update [Neumann, Sudholt, Witt'10]

All are **balanced!** (starting with the uniform distribution)

Compact Genetic Algorithm (cGA)

Univariate Marginal Distribution Algorithm (UMDA)

 $\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\boldsymbol{p}_{i}^{(t)} + \frac{1}{K}\left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)}\right) \mid \boldsymbol{p}^{(t)}\right]$

Compact Genetic Algorithm (cGA)

$$\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\boldsymbol{p}_{i}^{(t)} + \frac{1}{K}\left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)}\right) \mid \boldsymbol{p}^{(t)}\right]$$
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$$= \frac{1}{K^{2}} \cdot \operatorname{Var}\left[\boldsymbol{\hat{x}}_{i}^{(1)} - \boldsymbol{\hat{x}}_{i}^{(2)} \mid \boldsymbol{p}^{(t)}\right]$$

Compact Genetic Algorithm (cGA)

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$$= \frac{2}{K^{2}} \boldsymbol{p}_{i}^{(t)} \left(1 - \boldsymbol{p}_{i}^{(t)}\right)$$

Compact Genetic Algorithm (cGA)

Univariate Marginal Distribution Algorithm (UMDA)

$$\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\boldsymbol{p}_{i}^{(t)} + \frac{1}{K}\left(\widehat{\boldsymbol{x}}_{i}^{(1)} - \widehat{\boldsymbol{x}}_{i}^{(2)}\right) \mid \boldsymbol{p}^{(t)}\right]$$
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$$= \frac{2}{K^{2}} \boldsymbol{p}_{i}^{(t)} \left(1 - \boldsymbol{p}_{i}^{(t)}\right)$$

 p_i covers a distance of at most $\frac{1}{4}$ with constant probability within $\Theta(K^2)$ iterations

Compact Genetic Algorithm (cGA)

Univariate Marginal Distribution Algorithm (UMDA)

 $\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\frac{1}{\mu}\sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_{i}^{(k)} \mid \boldsymbol{p}^{(t)}\right]$

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p_i covers a distance of at most
$$\frac{1}{4}$$
 with constant probability within $\Theta(K^2)$ iterations

Compact Genetic Algorithm (cGA)

Univariate Marginal Distribution Algorithm (UMDA)

 $= \frac{1}{\mu^2} \operatorname{Var} \left[\sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_i^{(k)} \mid \boldsymbol{p}^{(t)} \right]$

 $\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\frac{1}{\mu}\sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_{i}^{(k)} \mid \boldsymbol{p}^{(t)}\right]$

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Univariate Marginal Distribution Algorithm (UMDA)

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 $= \frac{1}{\mu^2} \sum_{k \in [\mu]} \operatorname{Var} \left[\widehat{\boldsymbol{x}}_i^{(k)} | \boldsymbol{p}^{(t)} \right]$

 $\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\frac{1}{\mu}\sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_{i}^{(k)} \mid \boldsymbol{p}^{(t)}\right]$

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 p_i covers a distance of at most $\frac{1}{4}$ with constant probability within $\Theta(K^2)$ iterations

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 $= \frac{1}{\mu^2} \operatorname{Var} \left[\sum_{k \in [\mu]} \widehat{\boldsymbol{x}}_i^{(k)} \mid \boldsymbol{p}^{(t)} \right]$

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$$p_i$$
 covers a distance of at most $\frac{1}{4}$ with constant probability within $\Theta(K^2)$ iterations

$$=\frac{1}{\mu}\boldsymbol{p}_{i}^{(t)}\left(1-\boldsymbol{p}_{i}^{(t)}\right)$$
Commonly Studied EDAs (variances)

Compact Genetic Algorithm (cGA)

Univariate Marginal Distribution Algorithm (UMDA)

$$\operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\boldsymbol{p}_{i}^{(t)} + \frac{1}{K}\left(\hat{\boldsymbol{x}}_{i}^{(1)} - \hat{\boldsymbol{x}}_{i}^{(2)}\right) \mid \boldsymbol{p}^{(t)}\right] \qquad \operatorname{Var}\left[\boldsymbol{p}_{i}^{(t+1)} \mid \boldsymbol{p}^{(t)}\right] = \operatorname{Var}\left[\frac{1}{\mu}\sum_{k\in[\mu]}\hat{\boldsymbol{x}}_{i}^{(k)} \mid \boldsymbol{p}^{(t)}\right] \\ = \operatorname{Var}\left[\frac{1}{K}\left(\hat{\boldsymbol{x}}_{i}^{(1)} - \hat{\boldsymbol{x}}_{i}^{(2)}\right) \mid \boldsymbol{p}^{(t)}\right] \qquad = \frac{1}{\mu^{2}}\operatorname{Var}\left[\sum_{k\in[\mu]}\hat{\boldsymbol{x}}_{i}^{(k)} \mid \boldsymbol{p}^{(t)}\right] \\ = \frac{1}{K^{2}} \cdot \operatorname{Var}\left[\hat{\boldsymbol{x}}_{i}^{(1)} - \hat{\boldsymbol{x}}_{i}^{(2)} \mid \boldsymbol{p}^{(t)}\right] \qquad = \frac{1}{\mu^{2}}\sum_{k\in[\mu]}\operatorname{Var}\left[\hat{\boldsymbol{x}}_{i}^{(k)} \mid \boldsymbol{p}^{(t)}\right] \\ = \frac{1}{K^{2}} \cdot \left(\operatorname{Var}\left[\hat{\boldsymbol{x}}_{i}^{(1)} \mid \boldsymbol{p}^{(t)}\right] + \operatorname{Var}\left[\hat{\boldsymbol{x}}_{i}^{(2)} \mid \boldsymbol{p}^{(t)}\right]\right) \qquad = \frac{1}{\mu}\boldsymbol{p}_{i}^{(t)}\left(1 - \boldsymbol{p}_{i}^{(t)}\right) \\ = \frac{2}{K^{2}}\boldsymbol{p}_{i}^{(t)}\left(1 - \boldsymbol{p}_{i}^{(t)}\right)$$

 p_i covers a distance of at most $\frac{1}{4}$ with constant probability within $\Theta(K^2)$ iterations p_i covers a distance of at most $\frac{1}{4}$ with constant probability within $\Theta(\mu)$ iterations



» We wish genetic drift to be low for at least $\Theta(T)$ iterations



[Sudholt, Witt'16; Friedrich, Kötzing, K.'16; Doerr, Zheng'20]

- » We wish genetic drift to be low for at least $\Theta(T)$ iterations
- » p_i covers a distance of at most $\frac{1}{4}$ with probability at least $1 \exp\left(-\Theta\left(\frac{1}{\operatorname{Var} \cdot T}\right)\right)$ [Doerr, Zheng'20]



[Sudholt, Witt'16; Friedrich, Kötzing, K.'16; Doerr, Zheng'20]

- » We wish genetic drift to be low for at least $\Theta(T)$ iterations
- » p_i covers a distance of at most $\frac{1}{4}$ with probability at least $1 \exp\left(-\Theta\left(\frac{1}{\operatorname{Var} \cdot T}\right)\right)$ [Doerr, Zheng'20]
- **Choose** algorithm parameters such that $Var \in \Theta\left(\frac{1}{T \log n}\right)$ Subscripts Guarantee with high probability



[Sudholt, Witt'16; Friedrich, Kötzing, K.'16; Doerr, Zheng'20]

- » We wish genetic drift to be low for at least $\Theta(T)$ iterations
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- » The trick carries over to weak preferences of bit values [Doerr, Zheng'20]





 $\boldsymbol{x} \mapsto \sum_{i \in [n]} \boldsymbol{x}_i =: |\boldsymbol{x}|_1 \qquad \boldsymbol{x} \mapsto \max\{i \in [0..n] \mid \forall j \in [i]: \boldsymbol{x}_j = 1\}$



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Also: (non-exhaustive)

- » Analyses on noise [Friedrich, Kötzing, K., Witt'17; Lehre, Nguyen'19; Lehre, Nguyen'21; Kötzing, Radhakrishnan'22]
- » Analyses on deception [Lehre, Nguyen'19; Doerr, K.'21]
- » New EDAs [Doerr, K.'20; Ajimakin, Devi'23]
- » Multi-valued EDAs [Ben Jedidia, Doerr, K.'24; Adak, Witt'24]







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 $\mathbb{E}\left[\boldsymbol{p}_{i}^{(t+1)} - \boldsymbol{p}_{i}^{(t)} \mid \boldsymbol{p}^{(t)}\right] \geq \frac{2}{K} \boldsymbol{p}_{i}^{(t)} \left(1 - \boldsymbol{p}_{i}^{(t)}\right) \cdot \Pr\left[D = 0 \mid \boldsymbol{p}^{(t)}\right] \qquad \text{$`` lgnoring the case } D = 1 \text{ does not change the result asymptotically here}$

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The probability of each outcome in *I* is very roughly the same, that is, $\frac{1}{|I|} \approx \frac{1}{\sigma}$

» Note that
$$\sigma = \sqrt{\operatorname{Var}\left[X \mid \boldsymbol{p}^{(t)}\right]} =$$

$$\sqrt{\sum_{j\in[n]\setminus\{i\}} \boldsymbol{p}_j^{(t)} \left(1-\boldsymbol{p}_j^{(t)}\right)}$$

sampling variance

$$\Pr\left[D=0 \mid \boldsymbol{p}^{(t)}\right] = \Pr\left[X=Y \mid \boldsymbol{p}^{(t)}\right]$$

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Jensen's inequality:
$$\frac{\sum_{k \in L} a_k^2}{|L|} \ge \left(\frac{\sum_{k \in L} a_k}{|L|}\right)^2$$
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 $= E[\varphi^{(t)} - \varphi^{(t+1)} | \mathbf{p}^{(t)}] \gtrsim \frac{1}{K} \sqrt{\varphi^{(t)}} \qquad = E[T] \lesssim K \sqrt{n}$

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» For $K \in O\left(\frac{\sqrt{n}}{\log(n)\log\log n}\right)$ the run time is in $\Omega\left(K^{1/3}n + n\log n\right)$ [Lengler, Sudholt, Witt'18] » Covers high genetic drift 11.6

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cGA and UMDA on OneMax



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